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**MATH 135 Fall 2020: Extra Practice 5**
**Warm-Up Exercises**

**WE01.** Let  $\mathcal{U} = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$ ,  $A = \{2, 4, 6, 9\}$ , and  $B = \{4, 5, 6, 7\}$ .

(a) Calculate the following:

- i.  $A \cup B = \{2, 4, 5, 6, 7, 9\}$
- ii.  $A \cap B = \{4, 6\}$
- iii.  $\overline{A} = \mathcal{U} - A = \{1, 3, 5, 7, 8\}$
- iv.  $\overline{B} = \mathcal{U} - B = \{1, 2, 3, 8, 9\}$
- v.  $A - B = \{2, 9\}$
- vi.  $B - A = \{5, 7\}$

(b) Are  $A$  and  $B$  disjoint? *No, since 4 is in both sets.*

(c) Give a set  $C$  such that  $C \subseteq B$ . *Let  $C = B$ .*

(d) Give a set  $D$  such that  $D \subsetneq A$ . *Let  $D = \{2\}$ .*

**WE02.** Suppose  $S$  and  $T$  are two sets. Prove that if  $S \cap T = S$ , then  $S \subseteq T$ . Is the converse true?

*Proof.* Let  $S$  and  $T$  be arbitrary sets such that their intersection is  $S$ . We must show that any element of  $S$  is an element of  $T$ .

Consider an element  $s$  in  $S$ . But  $S$  is equal to  $S \cap T$ . Elements of an intersection are elements of the original sets, so  $s \in T$ , as desired.

For the converse, consider another two sets,  $S_1$  and  $T_1$ , where  $S_1 \subseteq T_1$ . This means that all elements of  $S_1$  are elements of  $T_1$ , that is, all elements of  $S_1$  are elements of both  $S_1$  and  $T_1$ . But this is just the definition of the intersection of  $S_1$  and  $T_1$ . Therefore, the converse is also true.  $\square$

**WE03.** Give an example of three sets  $A$ ,  $B$ , and  $C$  such that  $B \neq C$  and  $B - A = C - A$ .

*Solution.* Let  $A = \{1\}$ ,  $B = \{2\}$  and  $C = \{1, 2\}$ . Then,  $B - A = \{2\}$  and  $C - A = \{2\}$ .  $\square$

**Recommended Problems**

**RP01.** Let  $A$  be a subset of the universe  $\mathcal{U}$ . Prove that  $A \cup \overline{A} = \mathcal{U}$ .

*Proof.* Recall that the complement of a set  $\overline{S}$  with respect to a universe  $\mathcal{U}$  is defined as the set  $\{x \in \mathcal{U} : \neg(x \in S)\}$ . Recall also that the union of two sets  $X$  and  $Y$ , again with universe  $\mathcal{U}$ , is defined as the set  $X \cup Y = \{x \in \mathcal{U} : x \in X \vee x \in Y\}$ .

Then,  $A \cup \overline{A} = \{x \in \mathcal{U} : x \in A \vee \neg(x \in A)\}$ . The disjunction of any logical statement with its negation is a tautology, so this property is true for all elements of  $\mathcal{U}$ . Therefore, the resulting set is simply  $\mathcal{U}$ .  $\square$

**RP02.** Let  $S$  and  $T$  be two sets which are subsets of the universe  $\mathcal{U}$ . Prove that

$$(S \cup T) - (S \cap T) = (S - T) \cup (T - S).$$

*Proof.* Let  $S$  and  $T$  be arbitrary subsets of  $\mathcal{U}$ , and  $x$  be an arbitrary element of  $\mathcal{U}$  such that it is an element of the left-hand side. We prove by showing that the left and right-hand sides are subsets of another, that is, the following universally quantified biconditional holds:

$$\forall x \in \mathcal{U}, x \in (S \cup T) - (S \cap T) \iff x \in (S - T) \cup (T - S)$$

This can be done by rewriting both sides in set-builder notation and applying logical equivalencies.

$$\begin{aligned} (S \cup T) - (S \cap T) &= \{x \in \mathcal{U} : (x \in S \cup T) \wedge (x \notin S \cap T)\} \\ &= \{x \in \mathcal{U} : (x \in S \vee x \in T) \wedge \neg(x \in S \wedge x \in T)\} \\ &= \{x \in \mathcal{U} : (x \in S \vee x \in T) \wedge (\neg(x \in S) \vee \neg(x \in T))\} \end{aligned}$$

Now, distributing, we have the property:

$$(x \in S \wedge x \notin S) \vee (x \in S \wedge x \notin T) \vee (x \in T \wedge x \notin S) \vee (x \in T \wedge x \notin T)$$

which can be equivalently expressed by removing falsities:

$$(x \in S \wedge x \notin T) \vee (x \in T \wedge x \notin S).$$

Now, we can apply the definitions of unions and complements in reverse:

$$\begin{aligned} (S \cup T) - (S \cap T) &= \{x \in \mathcal{U} : (x \in S \wedge x \notin T) \vee (x \in T \wedge x \notin S)\} \\ &= \{x \in \mathcal{U} : (x \in S \wedge x \notin T)\} \cup \{x \in \mathcal{U} : (x \in T \wedge x \notin S)\} \\ &= (S - T) \cup (T - S) \quad \square \end{aligned}$$

**RP03.** Let  $A = \{n \in \mathbb{Z} : 2 \mid n\}$  and  $B = \{n \in \mathbb{Z} : 4 \mid n\}$ . Let  $n \in \mathbb{Z}$ . Prove that  $n \in (A - B)$  if and only if  $n = 2k$  for some odd integer  $k$ .

*Proof.* We prove the biconditional by proving both implications.

( $\Rightarrow$ ) Let  $n$  be an arbitrary integer element of  $A - B$ , i.e.,  $n \in A$  but  $n \notin B$ . Then, the defining property of  $A$  holds but that of  $B$  does not. Therefore,  $2 \mid n$  but  $4 \nmid n$ .

Since  $2 \mid n$ , it may be written as  $n = 2q$  for some integer  $q$ .

If  $q$  is even, then  $n = 2(2p)$  for some integer  $p$ . That means  $n = 4p$ , so  $n \mid 4$ , which is a contradiction. Therefore,  $q$  must be odd, and  $n$  may be written as  $n = 2k$  for an odd integer  $k = q$ .

( $\Leftarrow$ ) Let  $n$  be an arbitrary integer such that  $n = 2k$  for some odd integer  $k$ . It immediately follows that  $2 \mid n$  and  $n \in A$ .

Also, since  $k$  is odd,  $n = 2(2d + 1) = 4(d + \frac{1}{2})$  for another integer  $d$ .  $d + \frac{1}{2}$  will never be an integer, so  $4 \nmid n$ , which means  $n \notin B$ .

However,  $n \in A$  and  $n \notin B$  is precisely the definition of  $n \in (A - B)$ .

Therefore, since both implications hold, the statement is true. □

**RP04.** Prove that there exist sets  $A$ ,  $B$ , and  $C$  such that  $A \cup B = A \cup C$  and  $B \neq C$ .

*Proof.* Let  $A = \{1, 2\}$ ,  $B = \{1\}$ , and  $C = \{2\}$ . Clearly,  $B \neq C$ .

We have  $A \cup B = \{1, 2\}$  and  $A \cup C = \{1, 2\}$ , which are equal.  $\square$

**RP05.** Prove or disprove. If  $A$ ,  $B$ , and  $C$  are sets, then  $A \cap (B \cup C) = (A \cap B) \cup C$ .

*Solution.* Let  $A$ ,  $B$ , and  $C$  be arbitrary sets. We disprove by showing  $(A \cap B) \cup C$  is not a subset of  $A \cap (B \cup C)$ .

Let  $x$  be an element of  $C$  that is not an element of  $A$ . Then, it is clearly an element of  $(A \cap B) \cup C$ , since it is an element of  $C$  and all elements of either set in a union are elements of the union.

However, it is not an element of  $A \cap (B \cup C)$ . Since it is an intersection, all such elements are elements of  $A$ , which  $x$  is not.

Therefore,  $(A \cap B) \cup C \not\subseteq A \cap (B \cup C)$ . Set equality is defined by bidirectional subsets, so the sets cannot be equal.  $\square$

**RP06.** Prove there is a unique set  $T$  such that for every set  $S$ ,  $S \cup T = S$ .

*Proof.* We suppose that  $T = \emptyset$ , that is,  $T$  is the set with no elements, and prove it.

(Existence) Since there are no elements in  $T$ , it may be written as  $T = \{x : P\}$  where  $P$  is a false logical statement.

Now, the union  $S \cup T$  is  $\{x : x \in S \vee P\}$ . but a statement disjoined with false gives itself, so we have  $\{x : x \in S\}$ , which is just  $S$ .

(Uniqueness) Let  $A$  and  $B$  be empty sets.

Then,  $\forall x \in \mathcal{U}, x \in A \implies x \in B$  is vacuously true, since the hypothesis is always false by definition. Therefore,  $A \subseteq B$ .

Likewise,  $\forall x \in \mathcal{U}, x \in B \implies x \in A$  is vacuously true. Therefore,  $B \subseteq A$ .

Since both  $A$  and  $B$  are mutual subsets,  $A = B$ , and the empty set is unique.  $\square$

## Challenges

**C01.** The *symmetric difference* of two sets  $A$  and  $B$ , denoted  $A \triangle B$ , is defined as

$$A \triangle B = (A - B) \cup (B - A).$$

Prove that  $(A \triangle B) \triangle C = A \triangle (B \triangle C)$ .

*Proof.* We will prove using logical equivalences.

Consider the left-hand side. By the given definition,

$$\begin{aligned} (A \triangle B) \triangle C &= ((A - B) \cup (B - A)) \triangle C \\ &= (((A - B) \cup (B - A)) - C) \cup (C - ((A - B) \cup (B - A))) \end{aligned}$$

which is a mess, so we re-express as a logical expression in set-builder notation. That is,  $\{x : P(x)\}$  for some open sentence  $P(x)$ . For convenience, let  $a \equiv x \in A$ ,  $b \equiv x \in B$ , and  $c \equiv x \in C$ .

$$\begin{aligned} P(x) &\equiv (((a \wedge \neg b) \vee (b \wedge \neg a)) \wedge \neg c) \vee (c \wedge \neg((a \wedge \neg b) \vee (b \wedge \neg a))) \\ &\equiv (a \wedge \neg b \wedge \neg c) \vee (b \wedge \neg a \wedge \neg c) \vee (c \wedge \neg(a \wedge \neg b) \wedge \neg(b \wedge \neg a)) \\ &\equiv (a \wedge \neg b \wedge \neg c) \vee (b \wedge \neg a \wedge \neg c) \vee (c \wedge (\neg a \vee b) \wedge (\neg b \vee a)) \end{aligned}$$

We now digress from this (also enormous) expression to simplify the last term. Recall in Problem RP02, we proved  $(X \vee Y) \wedge (\neg X \vee \neg Y) \equiv (X \wedge Y) \vee (\neg X \wedge \neg Y)$ . We may now apply this equivalence with  $X \equiv \neg a$  and  $Y \equiv b$ .

$$\begin{aligned} P(x) &\equiv (a \wedge \neg b \wedge \neg c) \vee (b \wedge \neg a \wedge \neg c) \vee (c \wedge ((\neg a \wedge b) \vee (\neg b \wedge a))) \\ &\equiv (a \wedge \neg b \wedge \neg c) \vee (b \wedge \neg a \wedge \neg c) \vee (c \wedge \neg a \wedge b) \vee (c \wedge \neg b \wedge a) \\ &\equiv (a \wedge b \wedge c) \vee (a \wedge \neg b \wedge \neg c) \vee (\neg a \wedge b \wedge \neg c) \vee (\neg a \wedge \neg b \wedge c) \end{aligned}$$

Now, consider the right-hand side. By the given definition,

$$\begin{aligned} A \triangle (B \triangle C) &= A \triangle ((B - C) \cup (C - B)) \\ &= (A - ((B - C) \cup (C - B))) \cup (((B - C) \cup (C - B)) - A) \end{aligned}$$

which we may express as  $\{x : Q(x)\}$  for some open sentence  $Q(x)$ .

$$\begin{aligned} Q(x) &\equiv (a \wedge \neg((b \wedge \neg c) \vee (c \wedge \neg b))) \vee (((b \wedge \neg c) \vee (c \wedge \neg b)) \wedge \neg a) \\ &\equiv (a \wedge (\neg(b \wedge \neg c) \wedge \neg(c \wedge \neg b))) \vee ((b \wedge \neg c \wedge \neg a) \vee (c \wedge \neg b \wedge \neg a)) \\ &\equiv (a \wedge (\neg b \vee c) \wedge (\neg c \vee b)) \vee (\neg a \wedge b \wedge \neg c) \vee (\neg a \wedge \neg b \wedge c) \end{aligned}$$

Applying the identity we just discovered, namely,  $X \wedge (\neg Y \vee Z) \wedge (\neg Z \vee Y) \equiv (X \wedge Y \wedge Z) \vee (X \wedge \neg Y \wedge \neg Z)$ , for  $X \equiv a$ ,  $Y \equiv b$ , and  $Z \equiv c$ .

$$Q(x) \equiv (a \wedge b \wedge c) \vee (a \wedge \neg b \wedge \neg c) \vee (\neg a \wedge b \wedge \neg c) \vee (\neg a \wedge \neg b \wedge c)$$

but this is exactly  $P(x)$ . Therefore, the right-hand side may be expressed  $\{x : P(x)\}$ , which is precisely the left-hand side.  $\square$